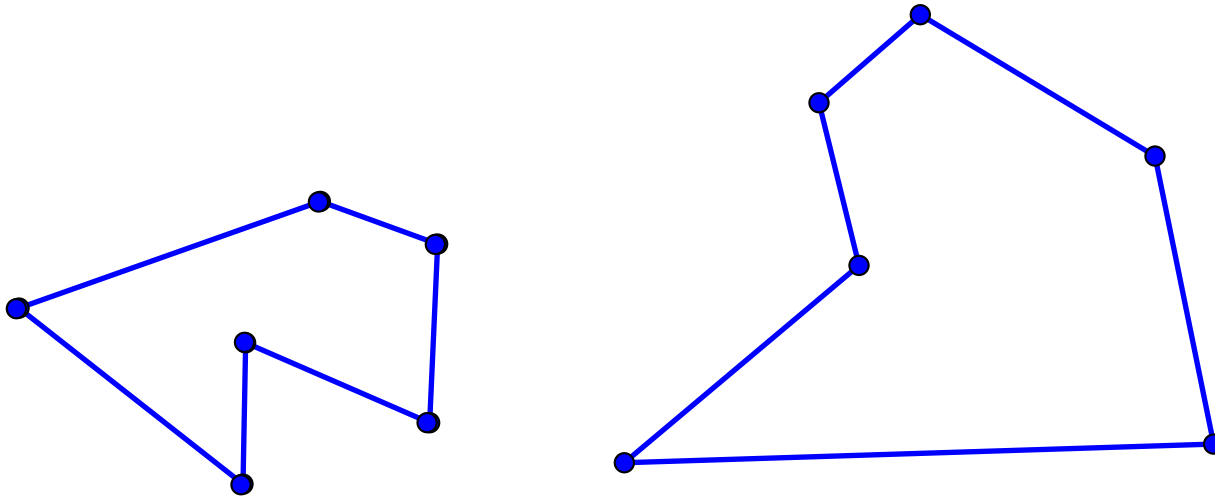


Points in Motion

Ileana Streinu
Smith College

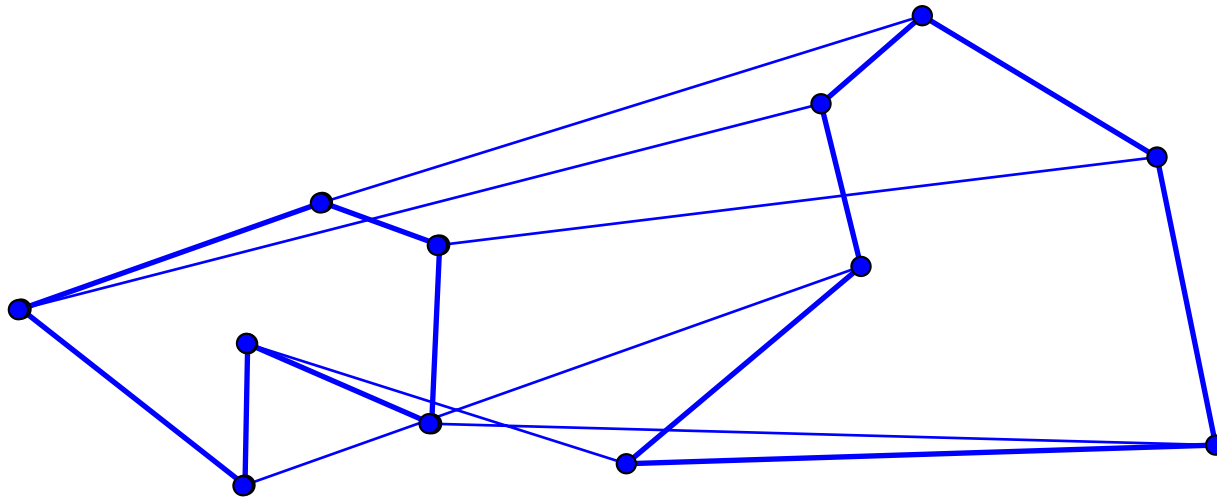
Motivation 1:

Morphing two polygonal shapes



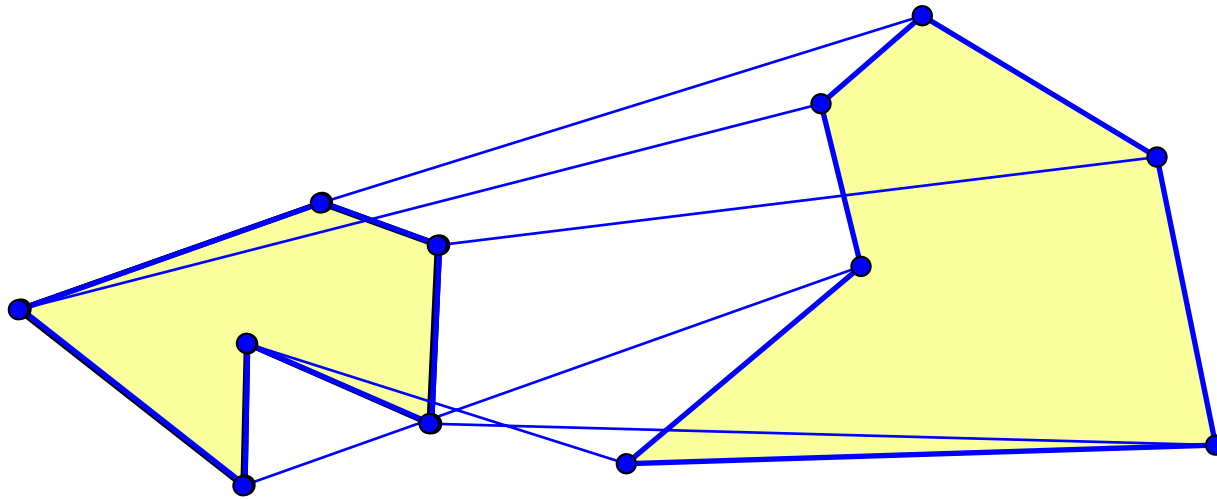
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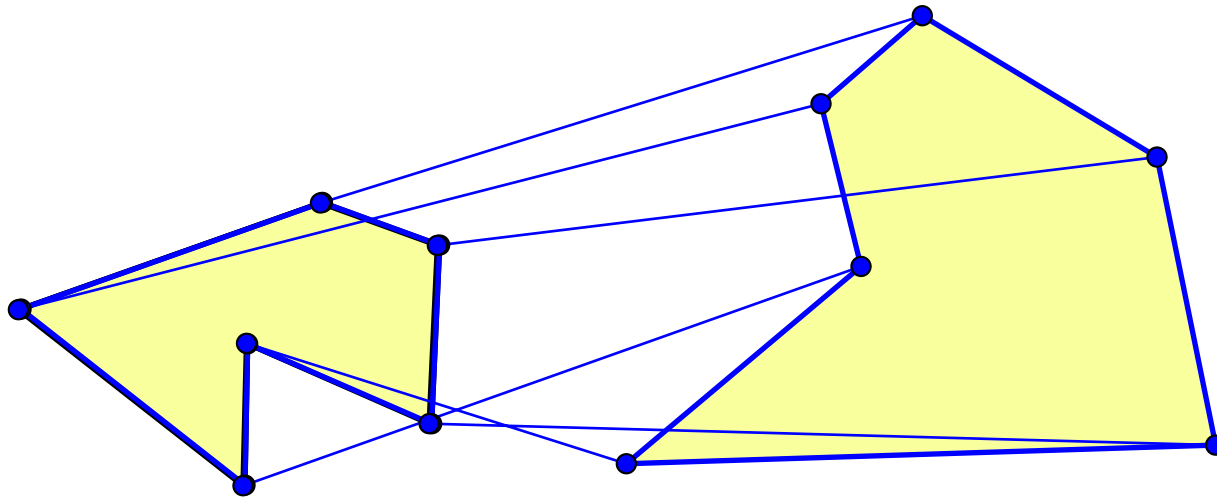
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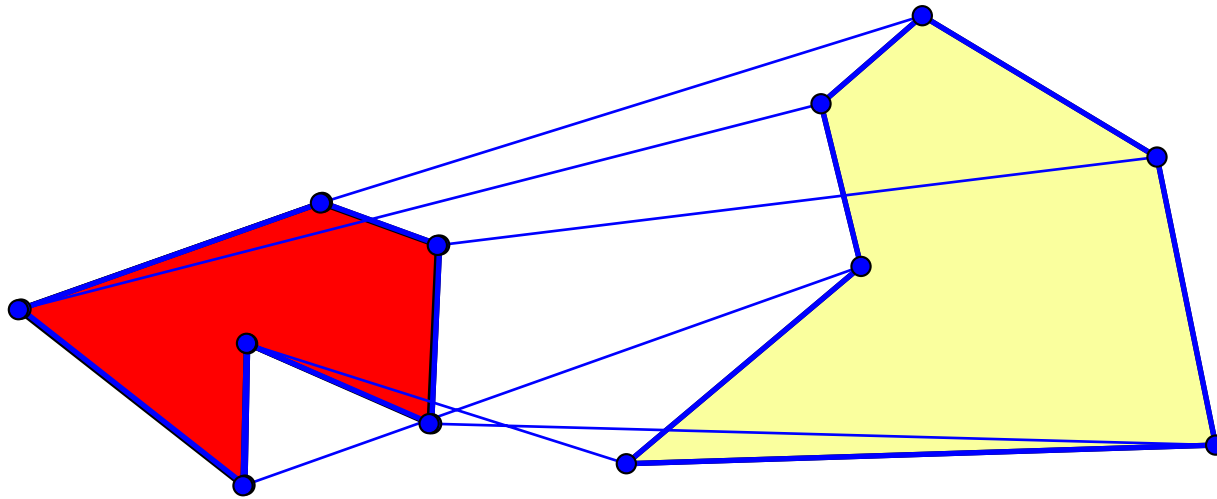
Morphing two polygonal shapes



Points move with constant velocities

Motivation 1:

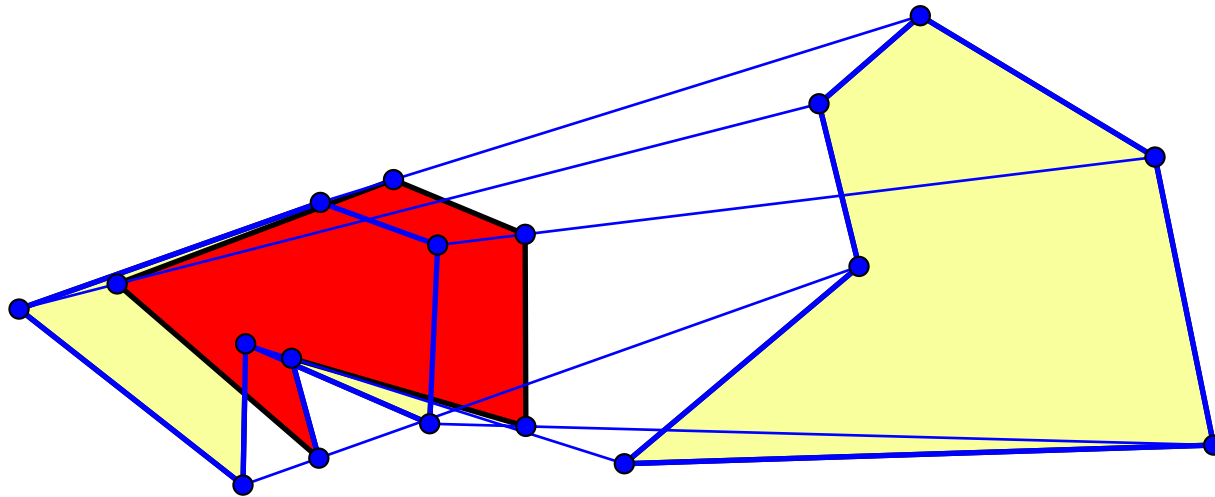
Morphing two polygonal shapes



Points move with constant velocities

Motivation 1:

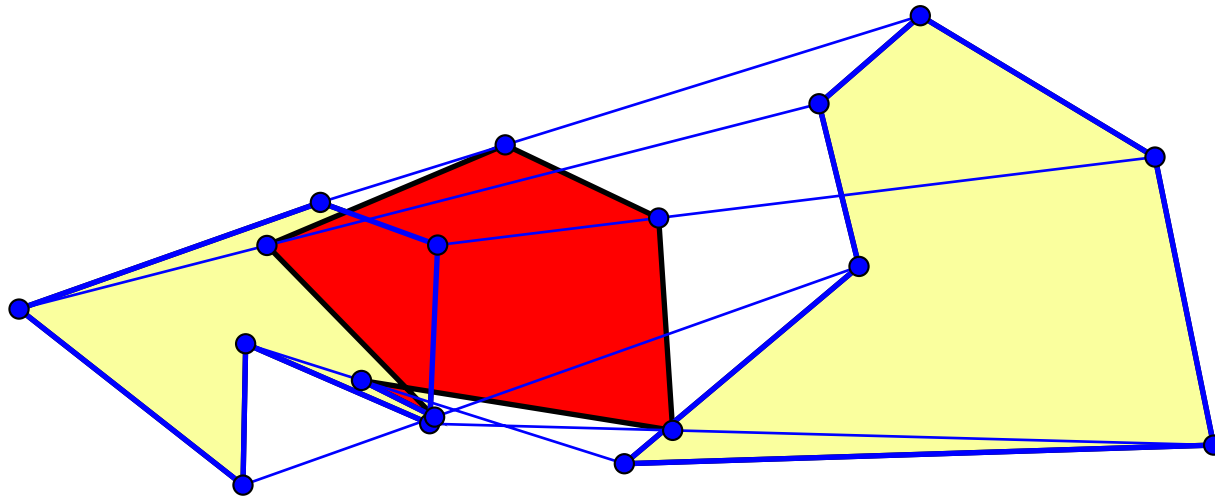
Morphing two polygonal shapes



Points move with constant velocities

Motivation 1:

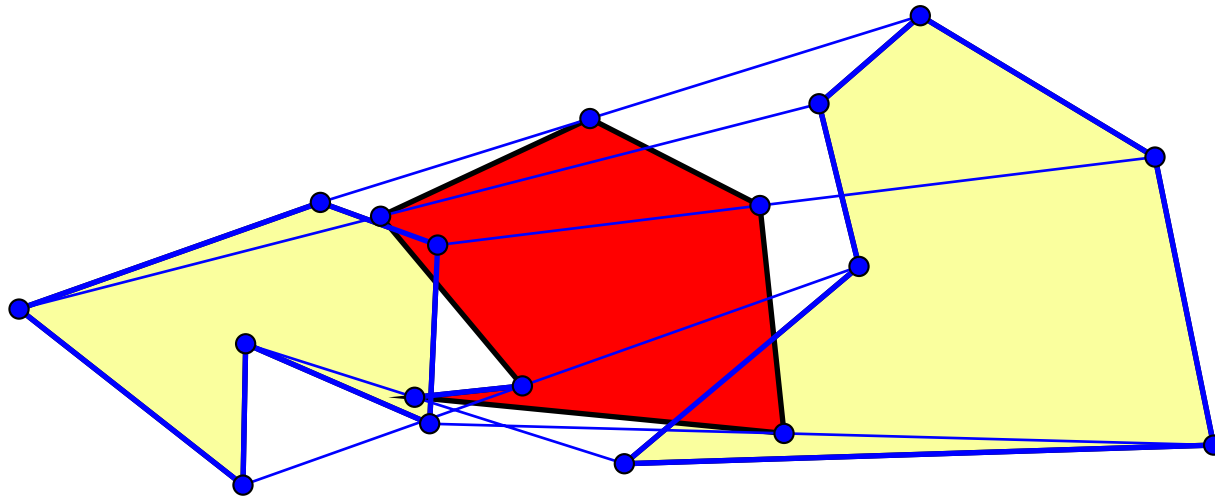
Morphing two polygonal shapes



Points move with constant velocities

Motivation 1:

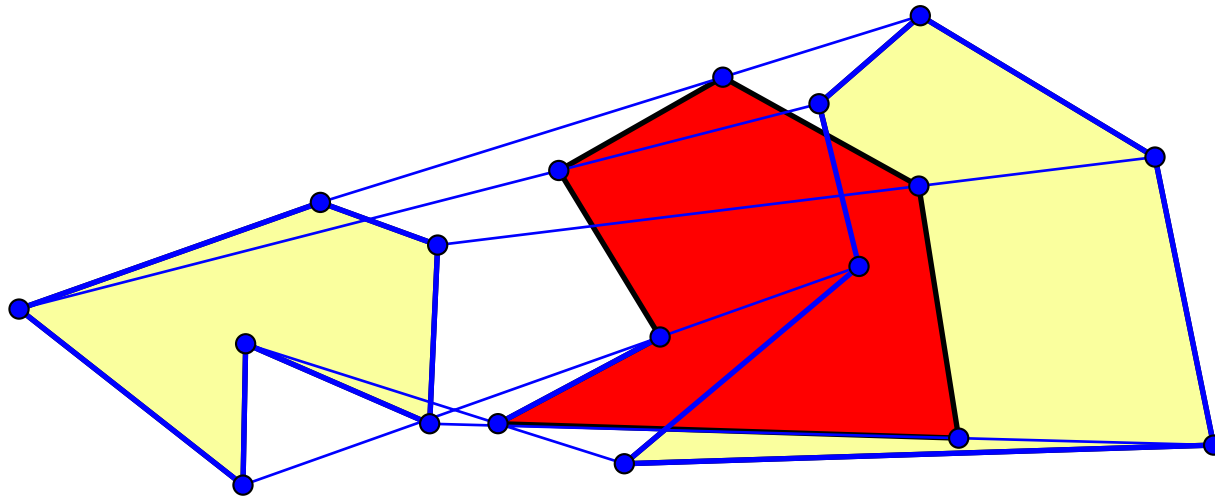
Morphing two polygonal shapes



Points move with constant velocities

Motivation 1:

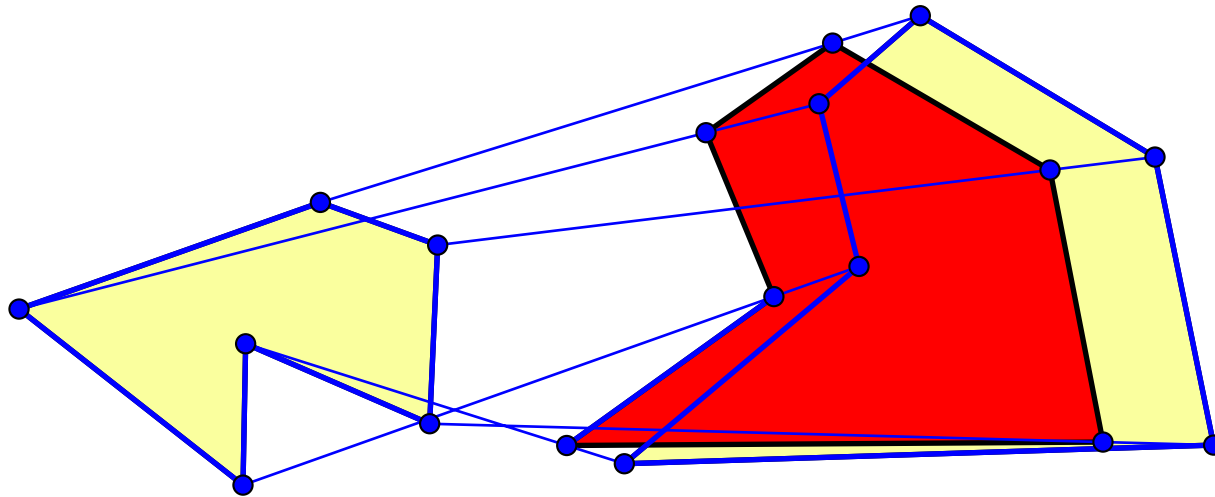
Morphing two polygonal shapes



Points move with constant velocities

Motivation 1:

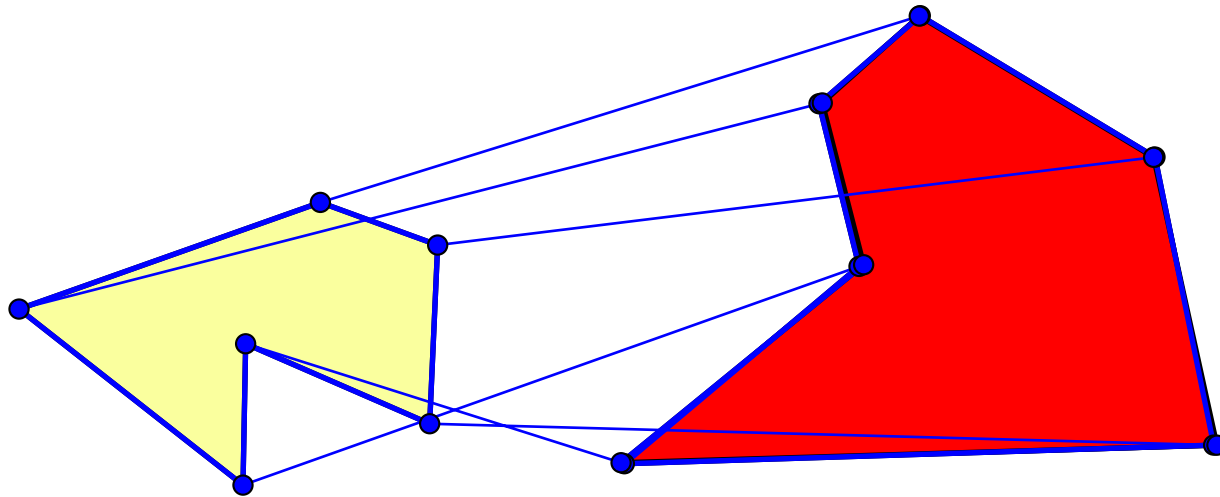
Morphing two polygonal shapes



Points move with constant velocities

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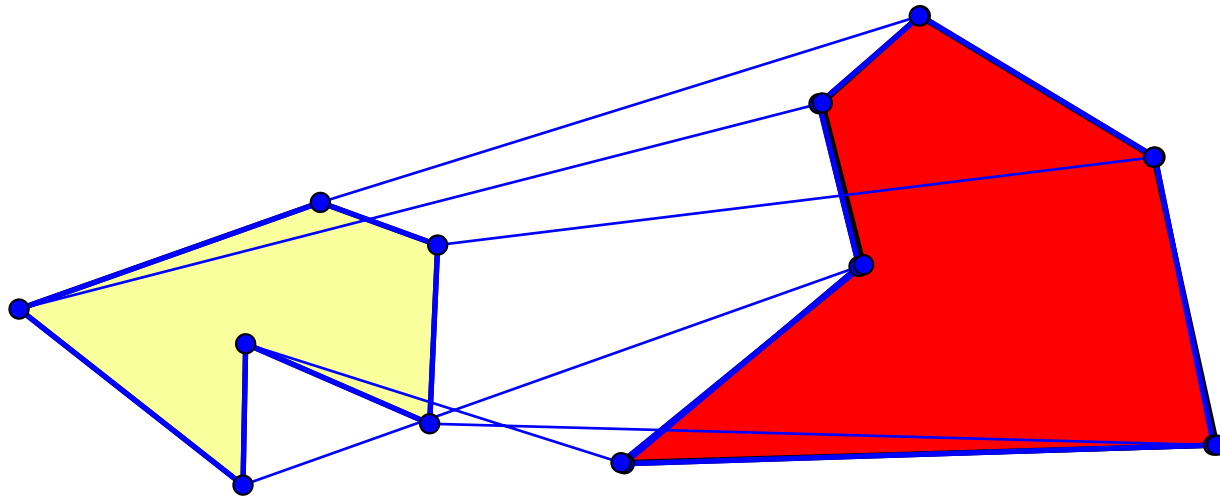
Morphing two polygonal shapes



Points move with constant velocities

Motivation 1:

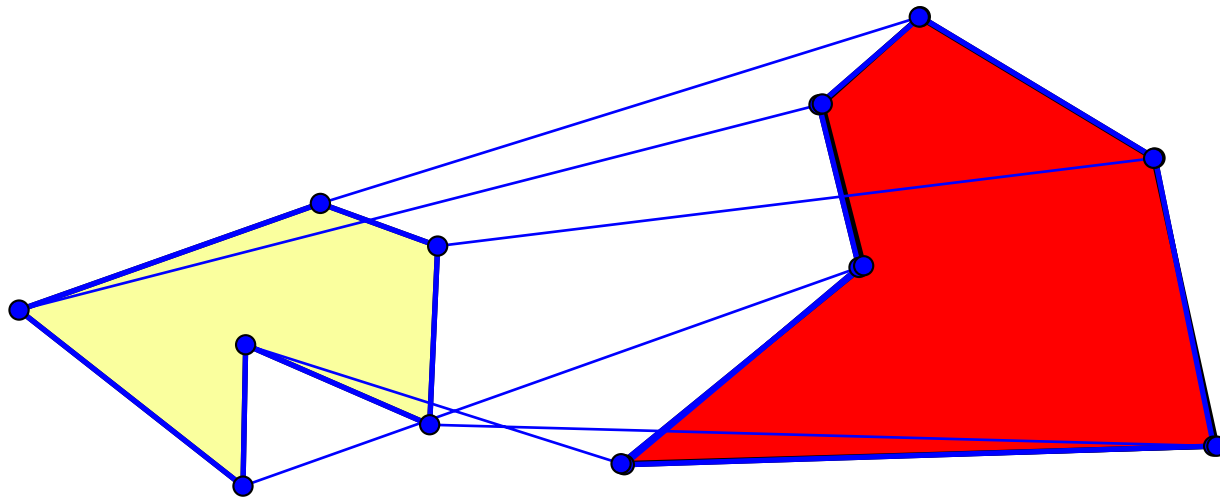
Morphing two polygonal shapes



Question: when can we GUARANTEE simplicity?

Motivation 1:

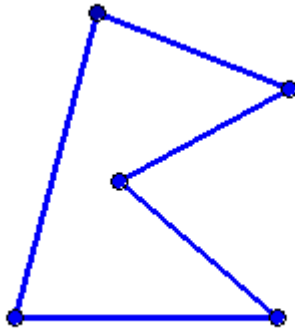
Morphing two polygonal shapes



Question: when can we **GUARANTEE** simplicity, for any duration of the constant-velocity motion?

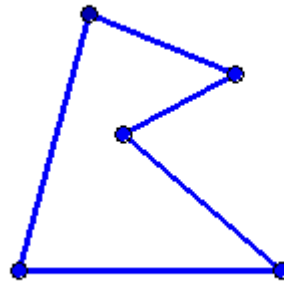
Motivation 2:

Morphing simple polygons with parallel edges (Guibas, Hershberger, Suri'00)



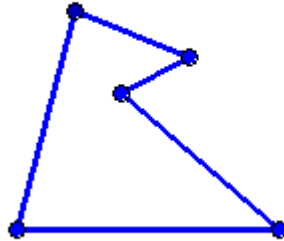
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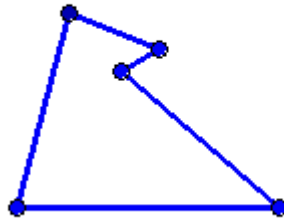
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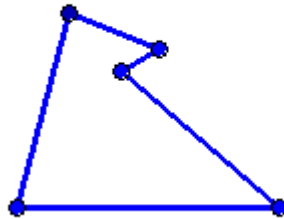
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Motivation 2:

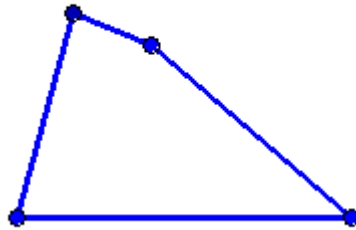
Morphing simple polygons with parallel edges (Guibas, Hershberger, Suri'00)



The motion is divided into pieces which are constant velocity, guaranteed to be collision free, and maintain the edge directions.

Motivation 2:

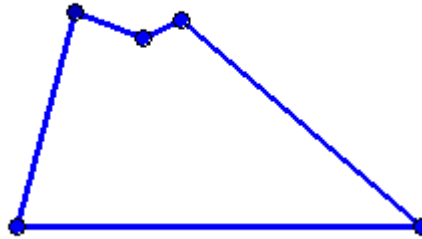
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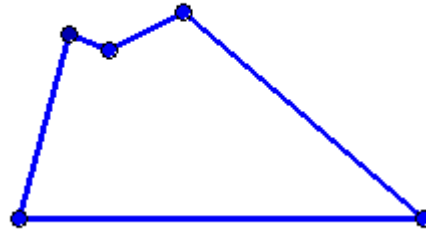
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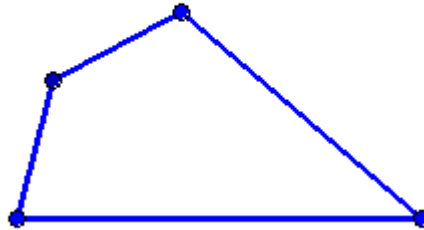
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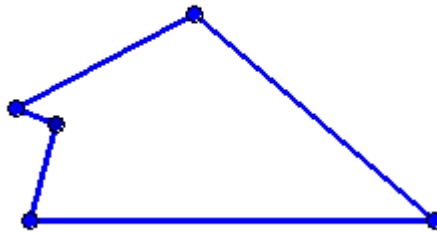
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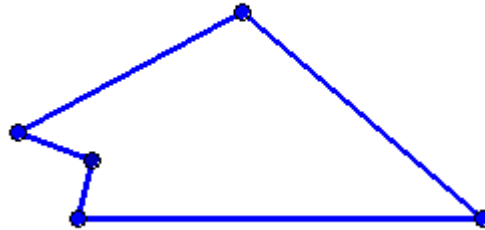
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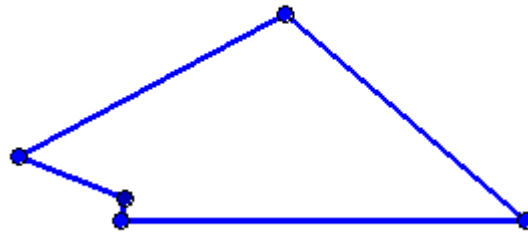
Morphing simple polygons with parallel edges (Guibas, Hershberger, Suri'00)



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Motivation 2:

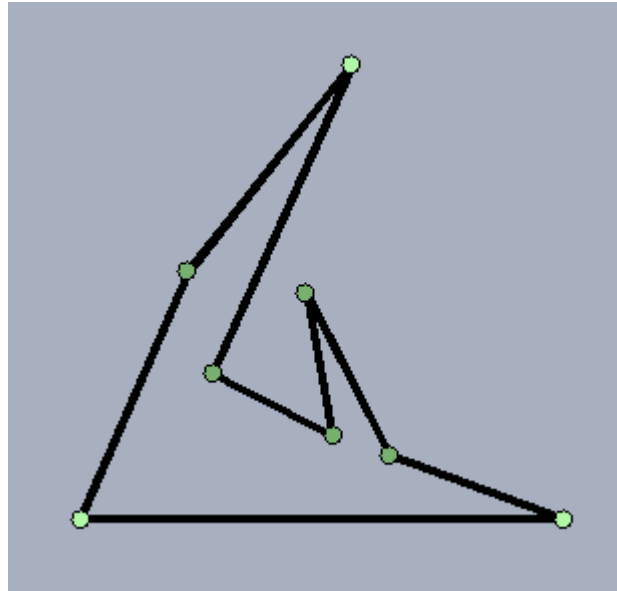
Morphing simple polygons with parallel edges (Guibas, Hershberger, Suri'00)



The motion is divided into pieces which are constant velocity, guaranteed to be collision free, and maintain the edge directions.

When can the whole morph be done with just **one** such motion?

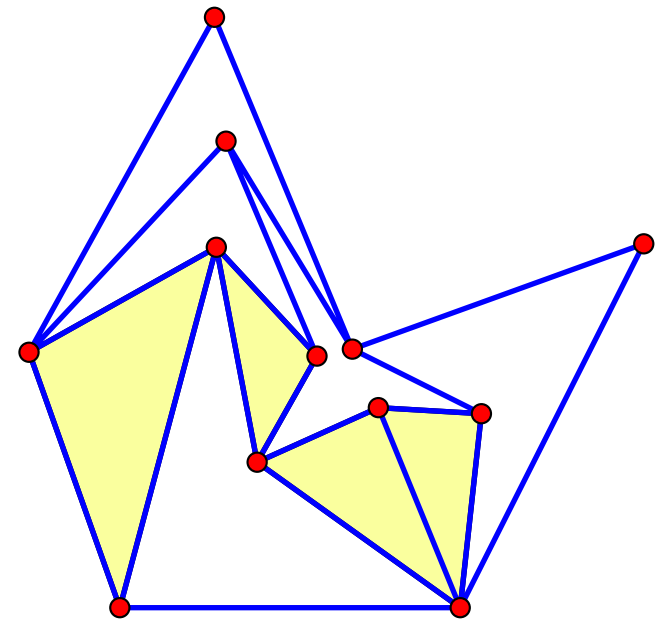
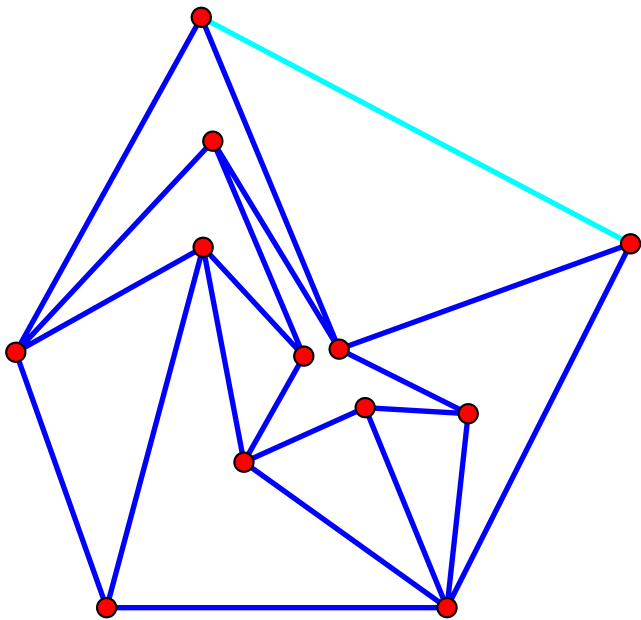
Example:



When can the whole morph be done with just one such motion?

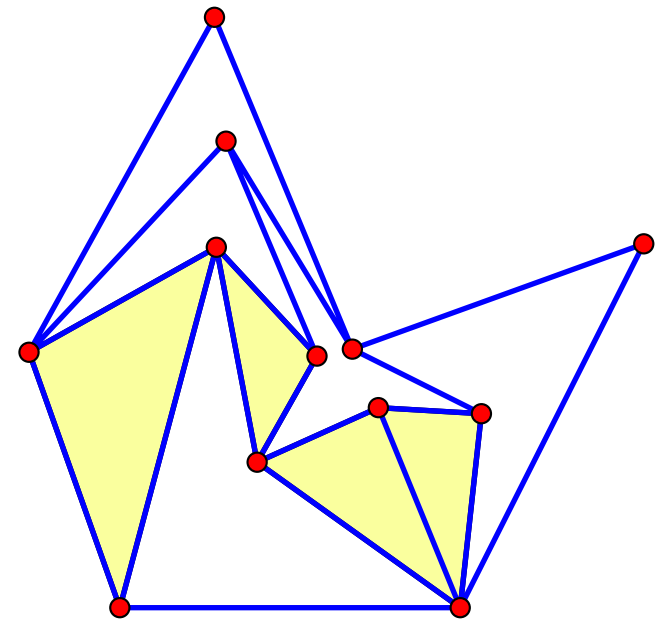
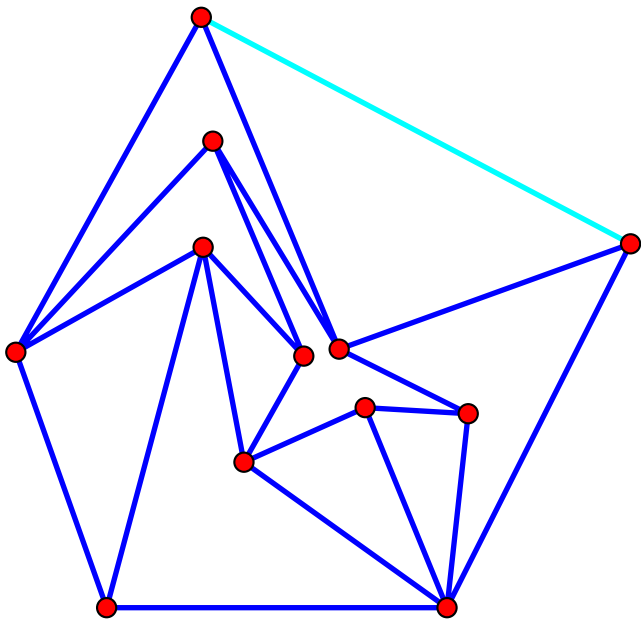
My real motivation ☺

Pointed pseudo-triangulation mechanisms (S'00)



My real motivation 😊

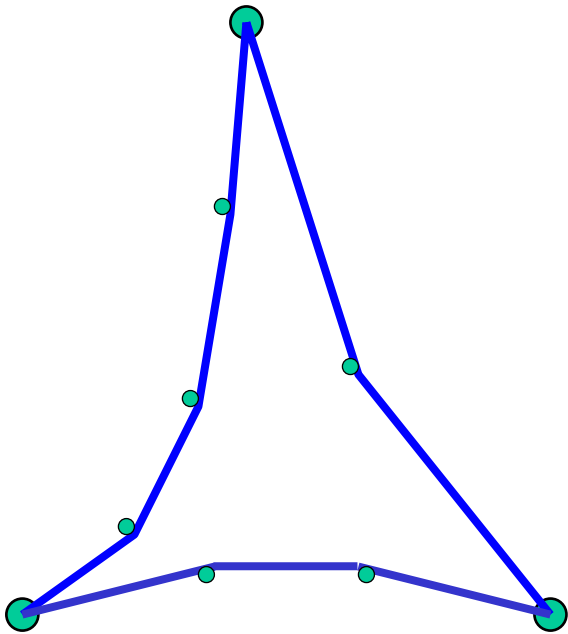
Pointed pseudo-triangulation mechanisms (S'00)



Pointed Pseudo-Triangulations: Definitions

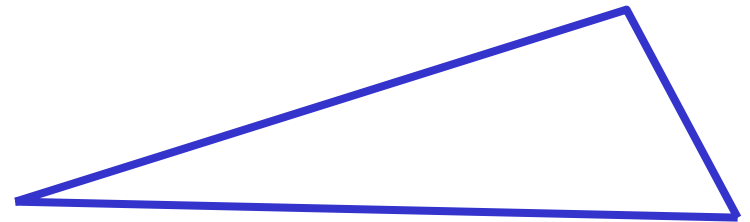
- Pseudo-Triangle
- Pointed Set of Edges
- PseudoTriangulation
- Pointed Pseudo-Triangulation

Pseudo Triangle



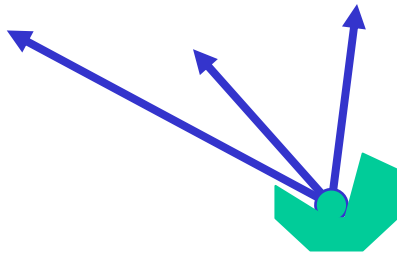
A simple polygon which has exactly **three** inner convex vertices.

In particular, a triangle is a pseudo-triangle.

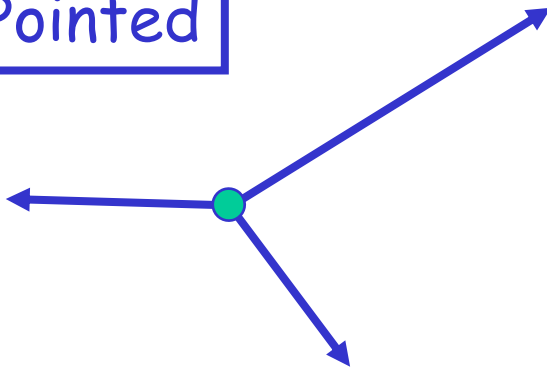


Pointed Planar Set of Vectors

Pointed

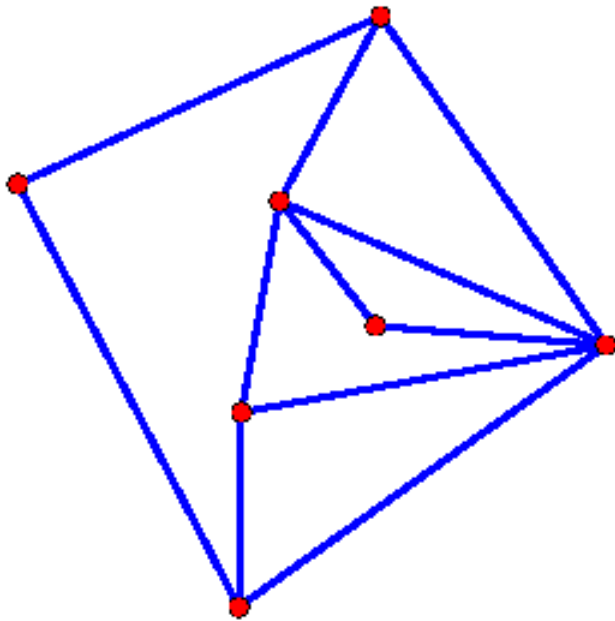


Not Pointed



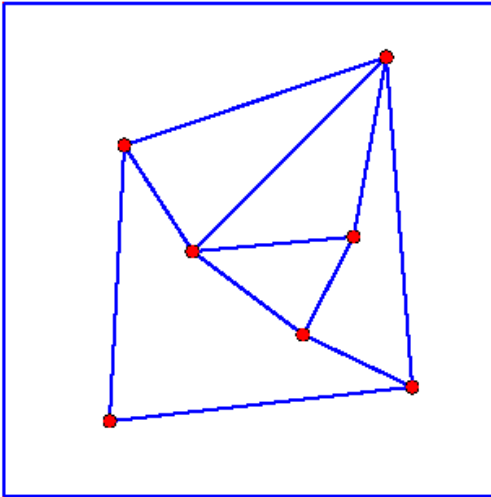
Two circularly adjacent vectors span a **reflex angle**

Pointed Pseudo Triangulation of a Planar Set of Points [S'00]

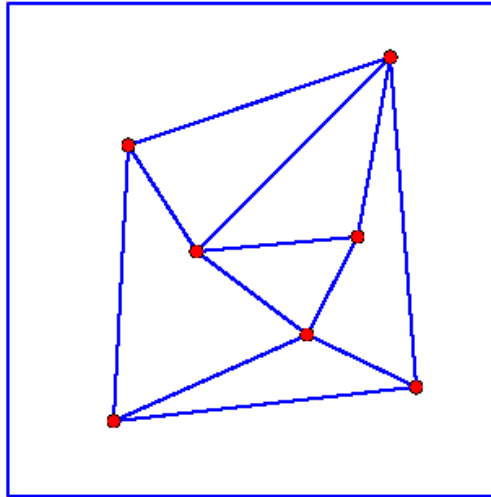


- Partitioning of the convex hull with a maximal set of non-crossing and pointed interior edges.
- The resulting faces are pseudo-triangles.

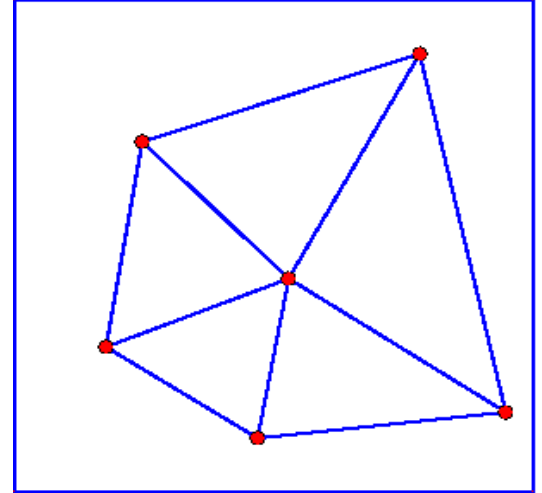
Other Pseudo Triangulations



Pointed



Not pointed



Not pointed

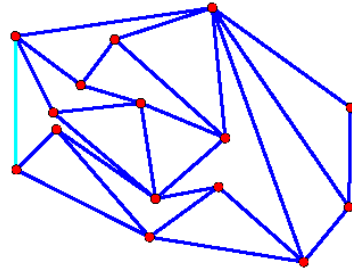
Pointed Pseudo Triangulations

Summary of main properties (S'00)

- Have exactly $2n-3$ edges, $n-2$ faces
- Are pointed, and **maximal** with the property of being both planar and pointed.
- Have the **hereditary Laman property**: any subset of k vertices is planar, pointed and has $\leq 2k-3$ edges
- Admit an inductive (**Henneberg**) construction.

Are **minimally rigid graphs**

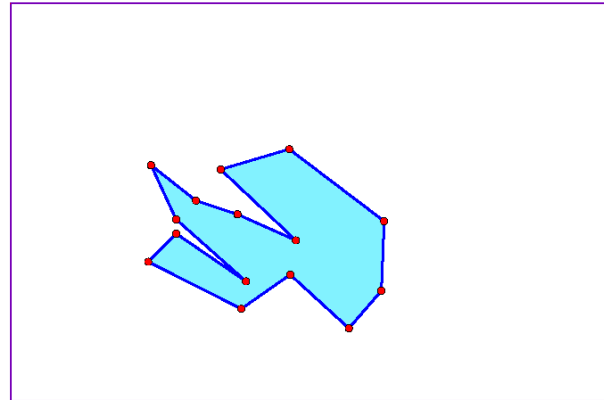
(with a special embedding)



Main **application** of Pointed Pseudo-Triangulations

A solution to the
Carpenter's Rule Problem

The Carpenter's Rule Problem



ODE-based:

Connelly, Demaine and Rote'00

See Erik Demaine's web page

<http://theory.lcs.mit.edu/~edemaine>

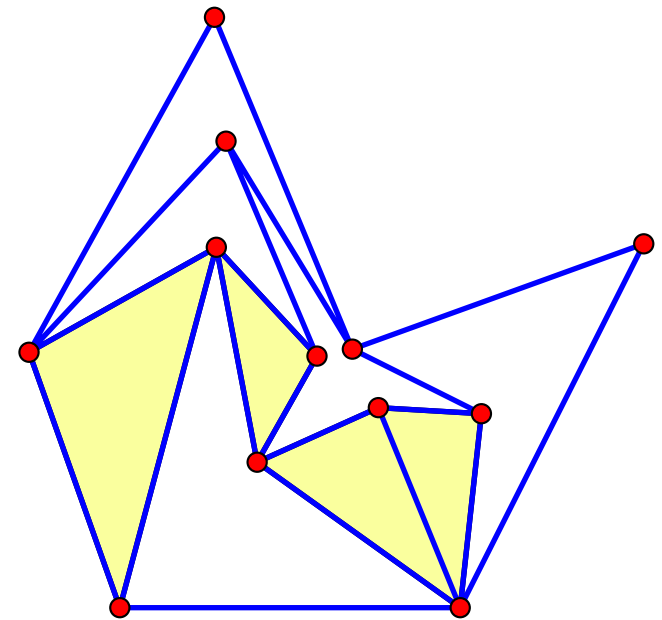
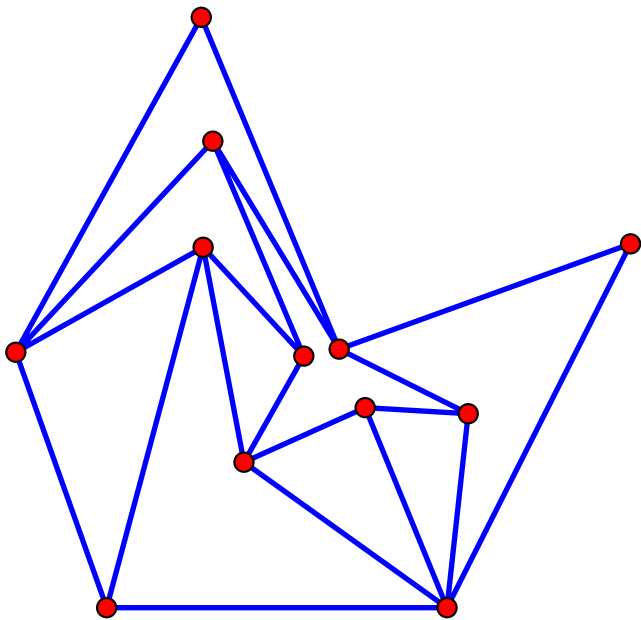
**Pseudo-triangulation based:
(S'00)**

From my web page

<http://cs.smith.edu/~streinu>








My real motivation 😊

Pointed pseudo-triangulation mechanisms (S'00)



As fixed-edge length mechanisms: expansive
As parallel redrawing mechanisms: WHAT?

Why study them?

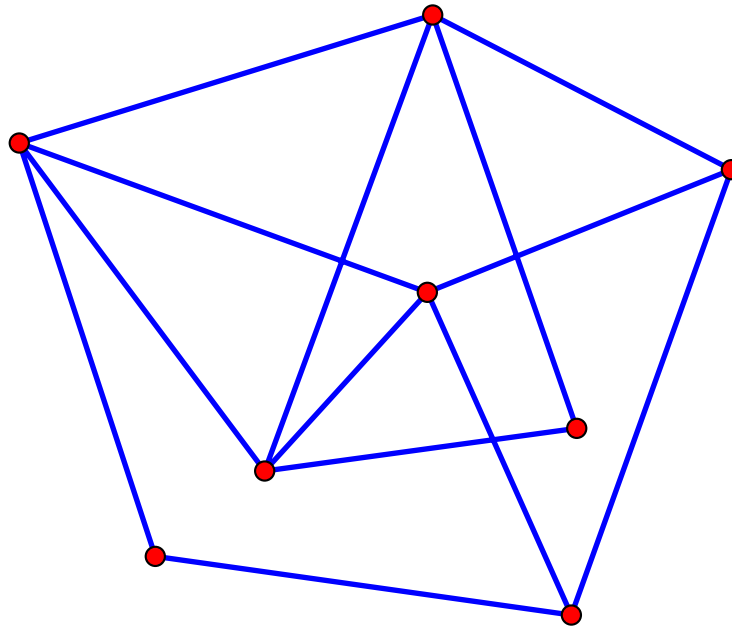
- Observed all ppt-mechanisms are planar (non-crossing) 
- Observed points move with constant velocities 
- Polygons on top of ppt mechanisms morph without proper crossings 
- Wanted to prove this 
- Ppt-mechanisms special case of 1dof Laman mechanism graphs: what about them? 
- Is kinetic planarity characterized by pointed pseudo triangulation mechanisms? 
- Apply it to other problems (morphing) 

Overview:

- The Parallel Redrawing model of rigidity:
 - fixed edge-direction, rather than
 - fixed edge-length
- Objects of study:
 - 1dof Laman mechanisms
 - Pointed pseudo-triangulation mechanisms
- Restricted to **GENERIC** situations
- Kinetic objects
 - Points
 - Embedded graphs
 - Polygons
- Focusing on:
 - Collisions
 - Edge crossings
 - Combinatorial invariants:
 - Rigid components
 - Oriented matroidal invariants:
 - » Partial hyperline sequences
 - » Combinatorial Pseudo-triangulations
- Algorithms:
 - Parallel redrawing sweep
- Further directions:
 - Kinetic point sets
 - Kinetic graphs
 - Combinatorial parallel redrawing sweep for combinatorial pseudo-triangulation mechanisms

Background:

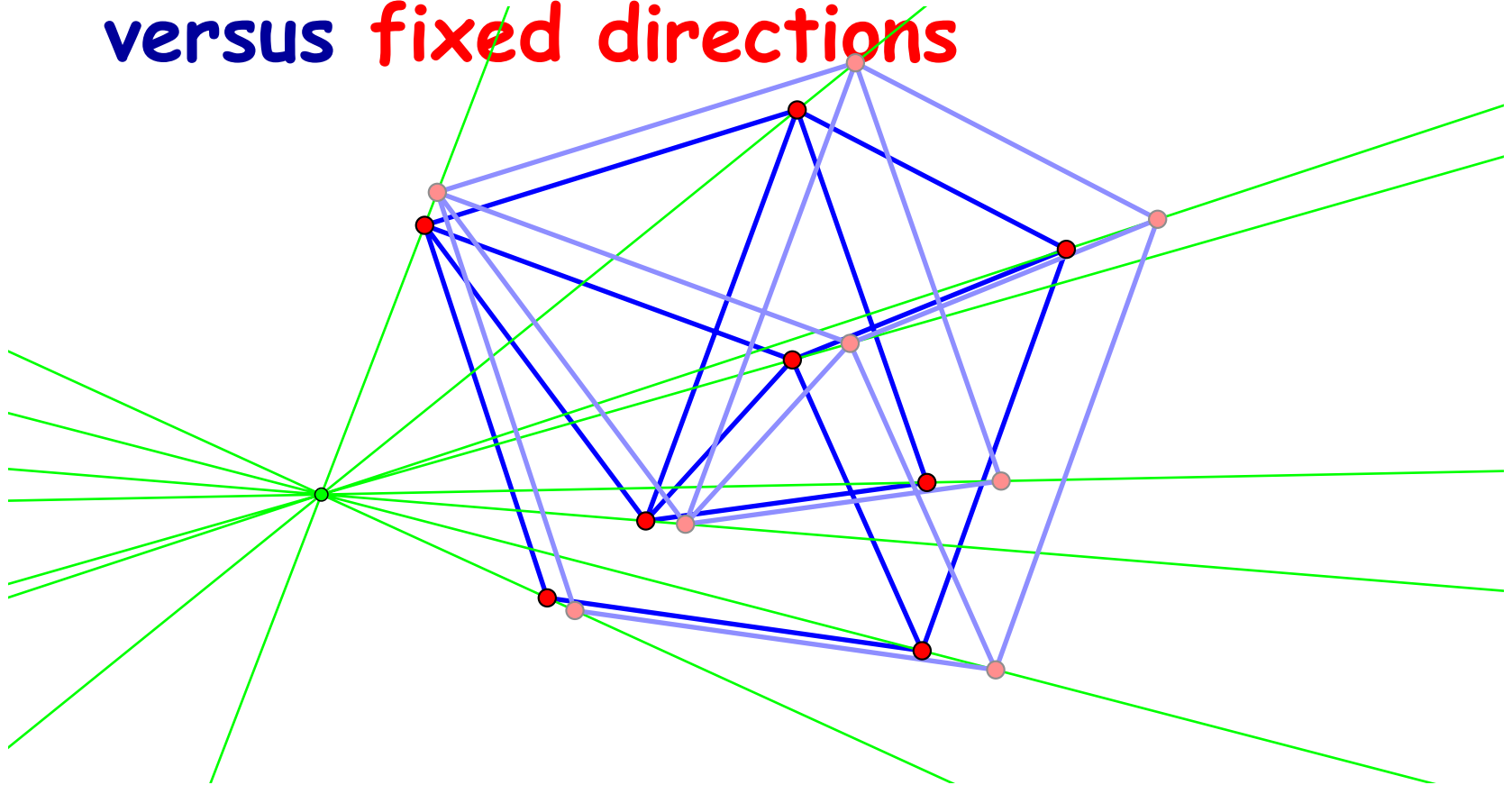
Rigidity with fixed **edge-lengths**
versus fixed **directions**



Minimally Rigid (Laman) graph: $2n-3$ edges,
every k -subset spans $\leq 2k-3$ edges

Background:

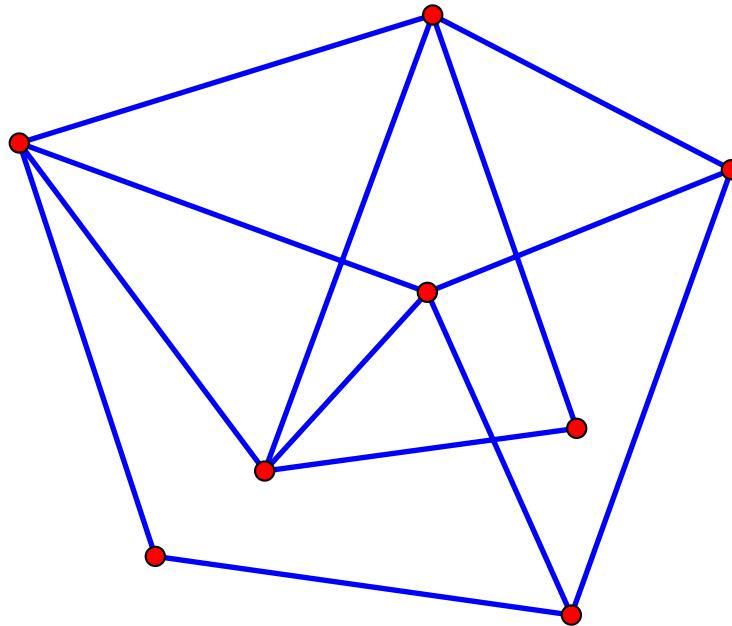
Rigidity with fixed edge-lengths
versus **fixed directions**



Laman graphs have only trivial parallel redrawings

Background:

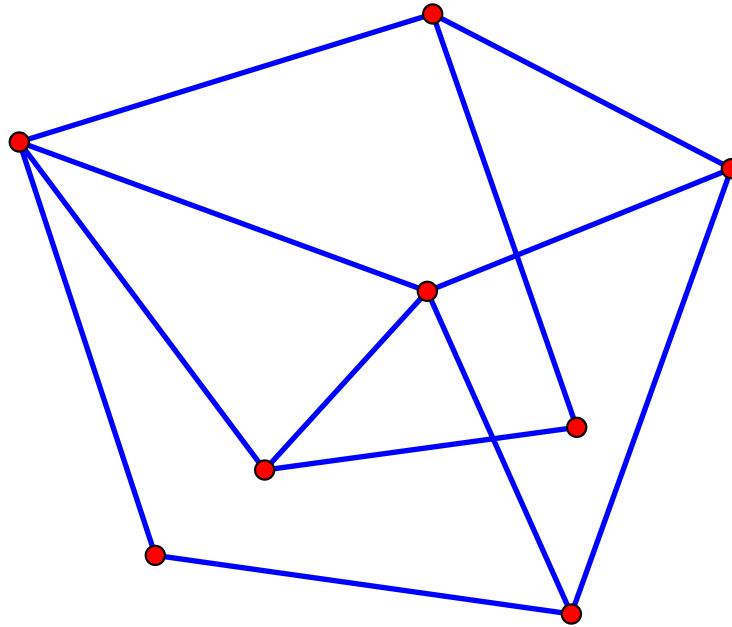
**Rigidity with fixed *edge-lengths*
versus fixed directions**



Minimally Rigid

Background:

Rigidity with fixed **edge-lengths**
versus fixed directions

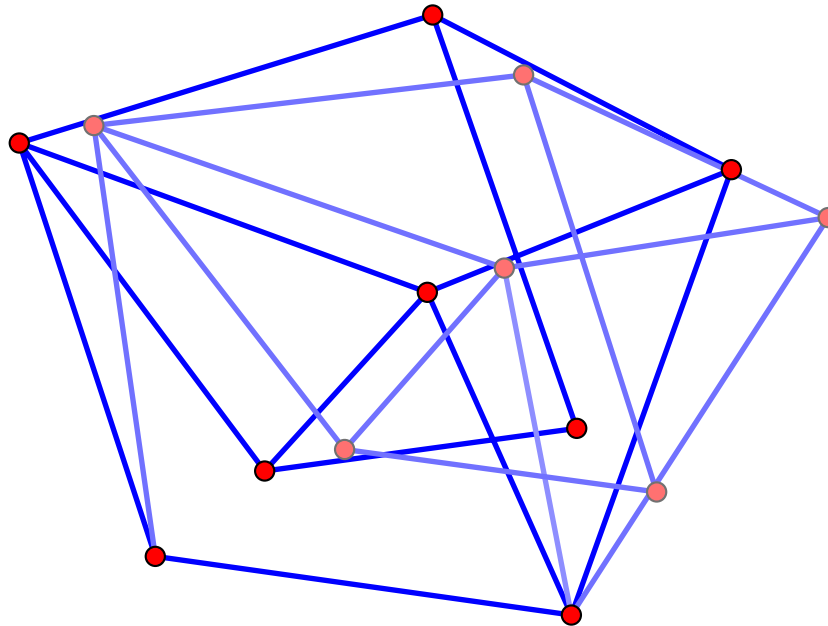


Laman mechanism

1dof flexible

Background:

Rigidity with fixed **edge-lengths
versus fixed directions**

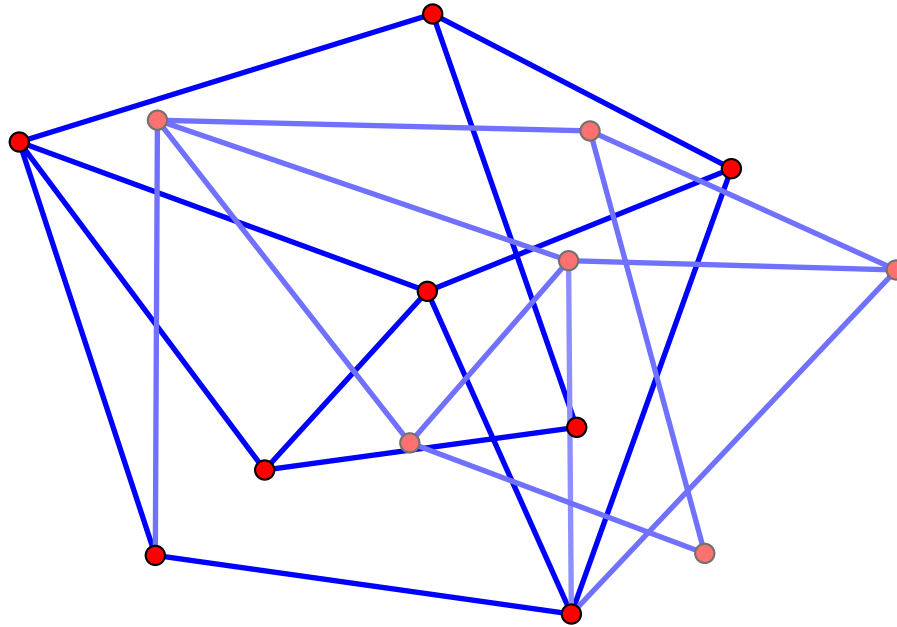


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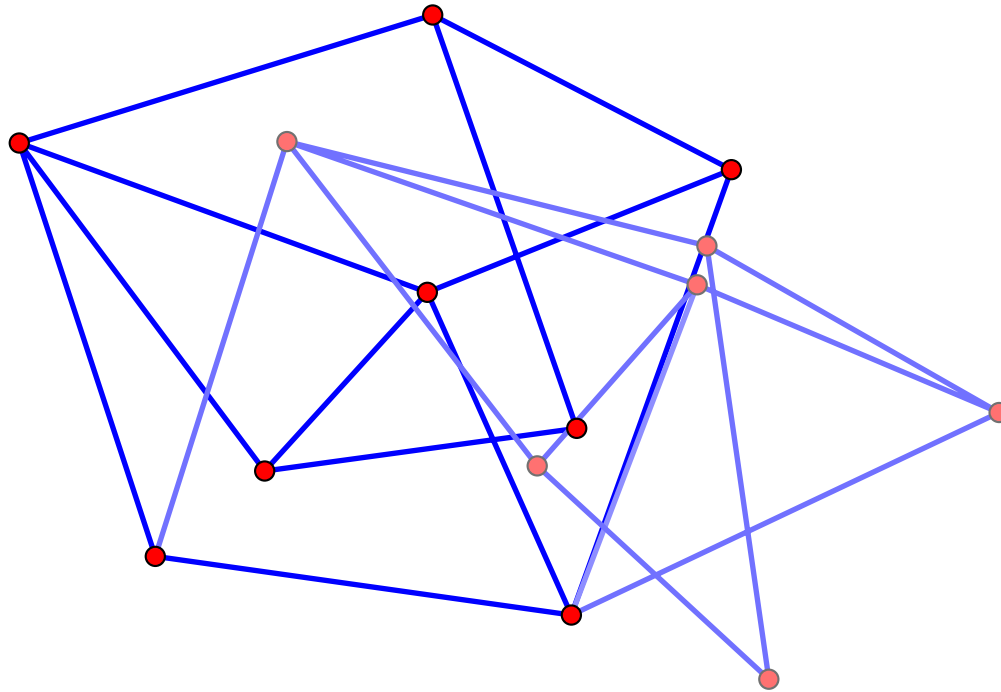


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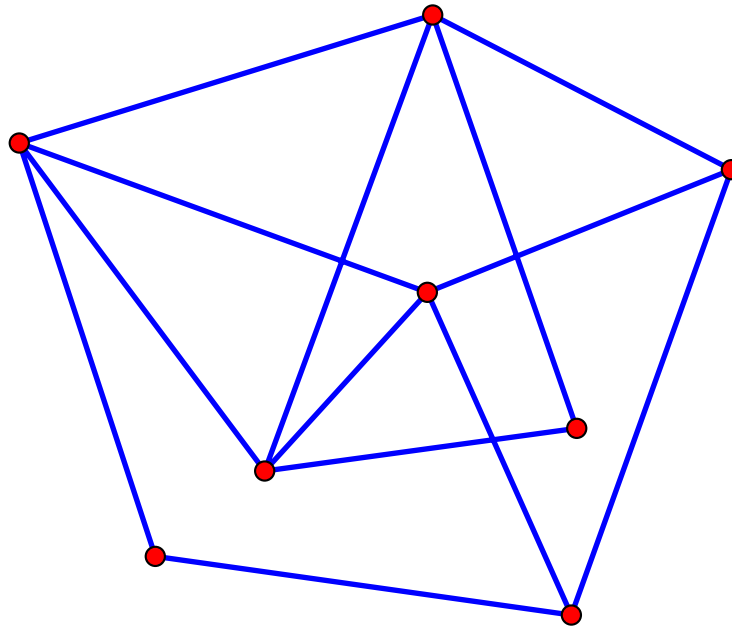


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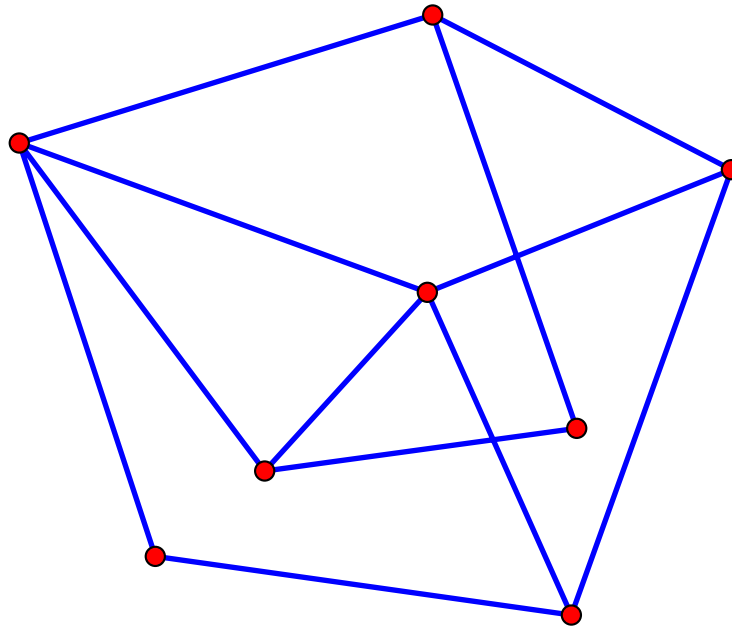


Laman graph:

only trivial parallel redrawings

Background:

**Rigidity with fixed edge-lengths
versus fixed directions**

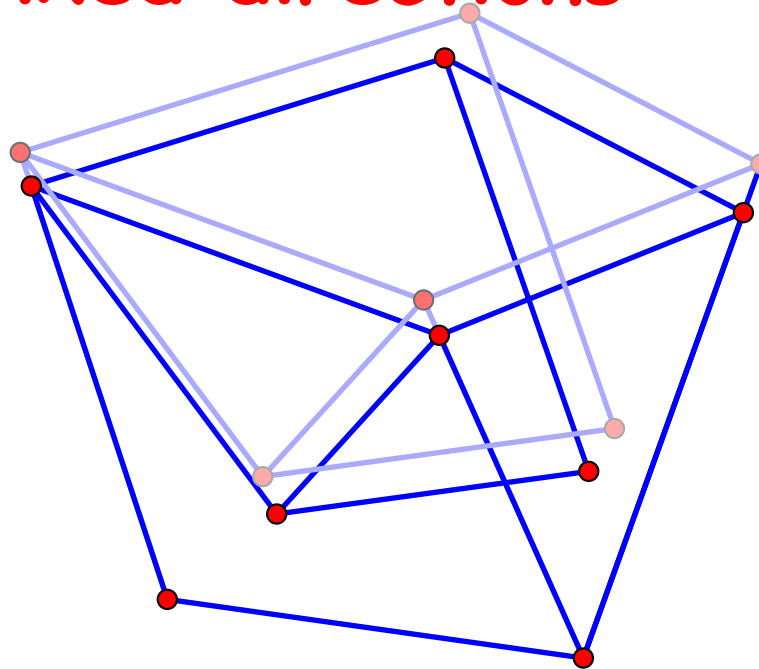


Laman mechanism (1dof flexible)

Non-trivial parallel redrawing

Background:

**Rigidity with fixed edge-lengths
versus fixed directions**

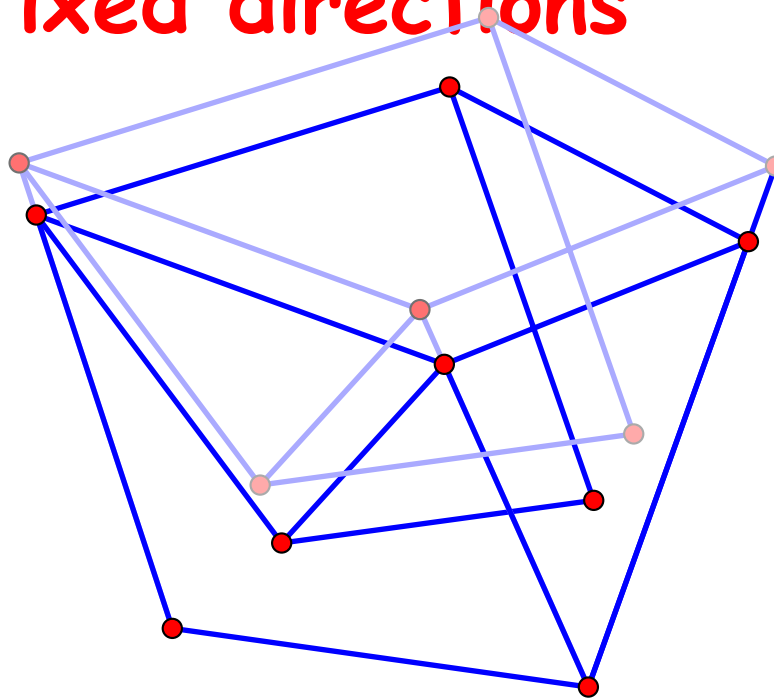


Laman mechanism (1dof flexible)

Non-trivial parallel redrawing

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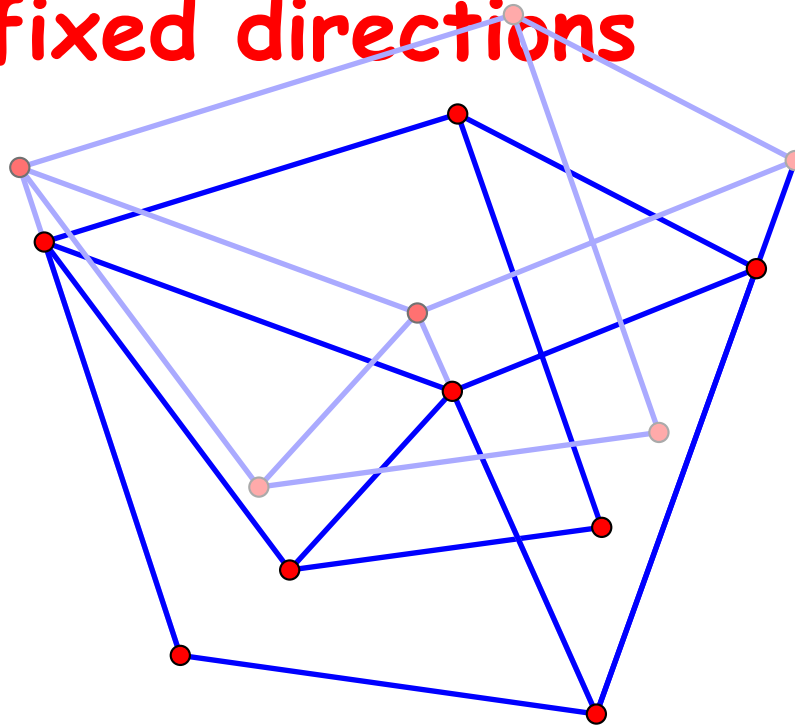


Laman mechanism (1dof flexible)

Non-trivial parallel redrawing

Background:

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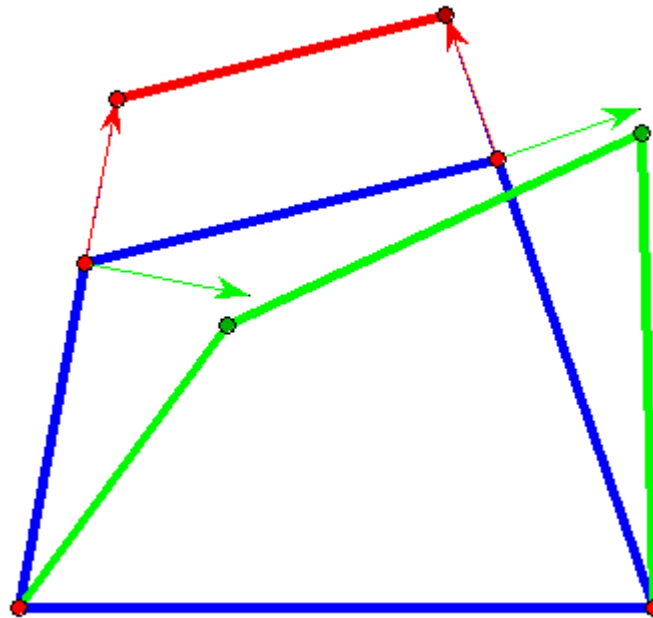


Laman mechanism (1dof flexible)

Non-trivial parallel redrawing

Background:

Relationship between motions in the two models



Orthogonal to each other

Reference: "folklore" from 19th century,
see Whiteley, Matroids survey

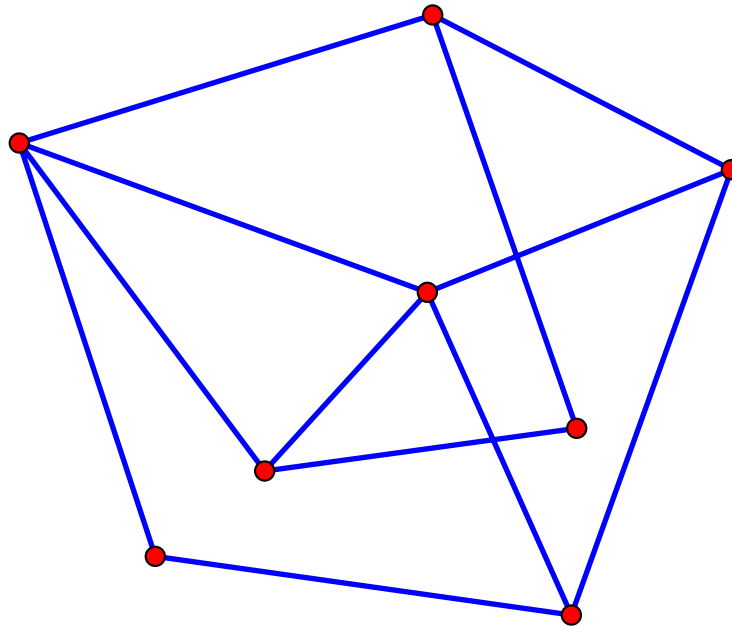
Plan

- **Configuration spaces** of parallel redrawing Laman graphs and 1dof Laman mechanisms
- **Parallel Redrawing Sweep** for 1dof Laman mechanisms
- **Pointed pseudo-triangulation** mechanisms
- Further problems: **kinetic point sets** and **graphs, combinatorial sweeps**

Plan

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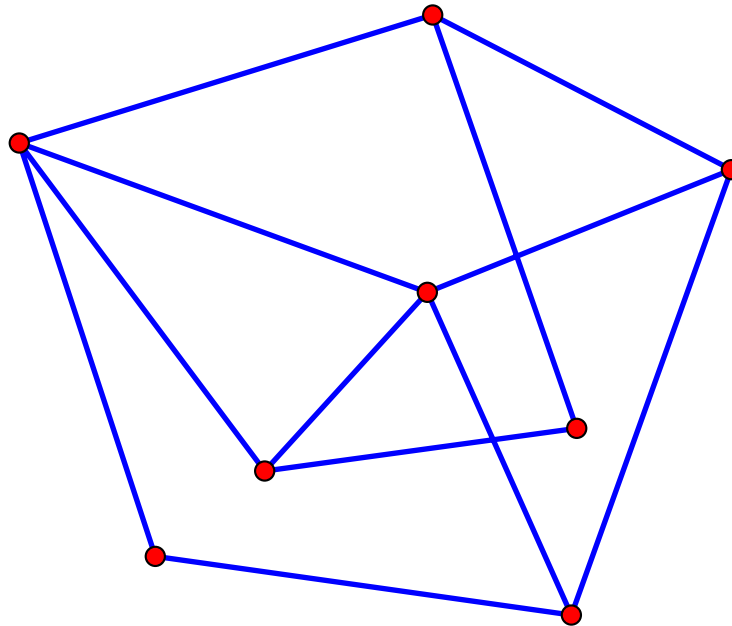
Direction networks and their parallel redrawings



Direction network (G, D) :

- Graph G
- Set D of directions (slopes) for the edges

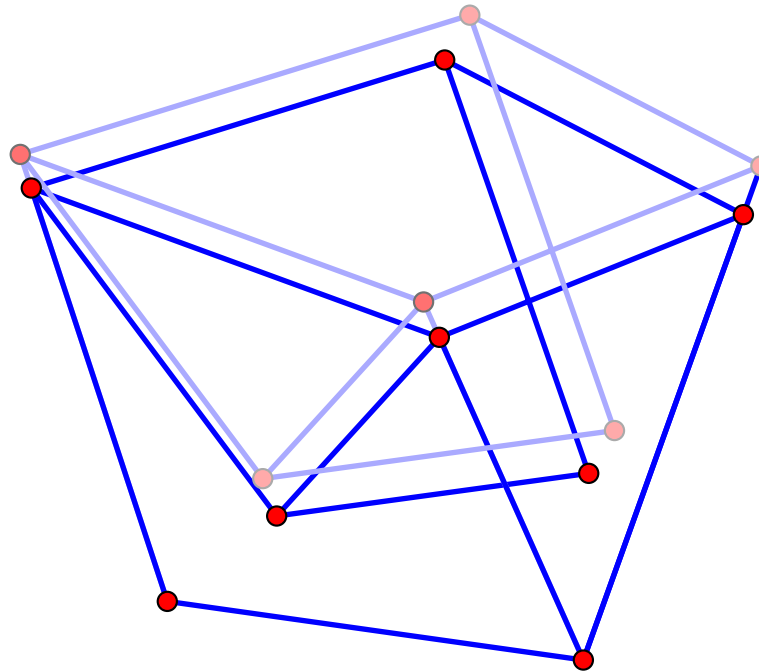
Direction networks and their parallel redrawings



Realization (embedding) of a direction network

- Mapping of vertices to points, edges to segments
- Consistent with given directions

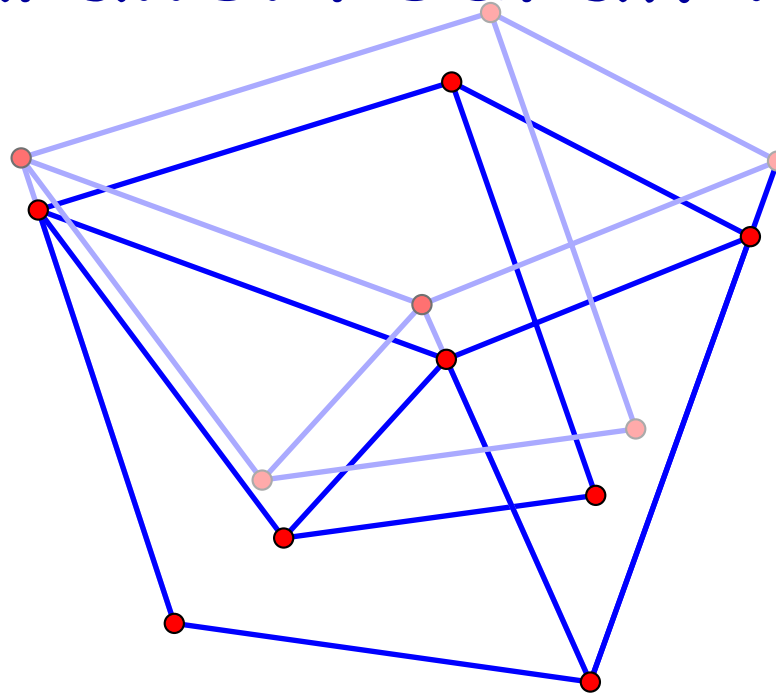
Direction networks and their parallel redrawings



Parallel redrawing of an embedded graph:

- Another realization of its underlying direction network
- Always can obtain similar ones by translation and rescaling
- Interesting: non-similar (non-trivial) parallel redrawings

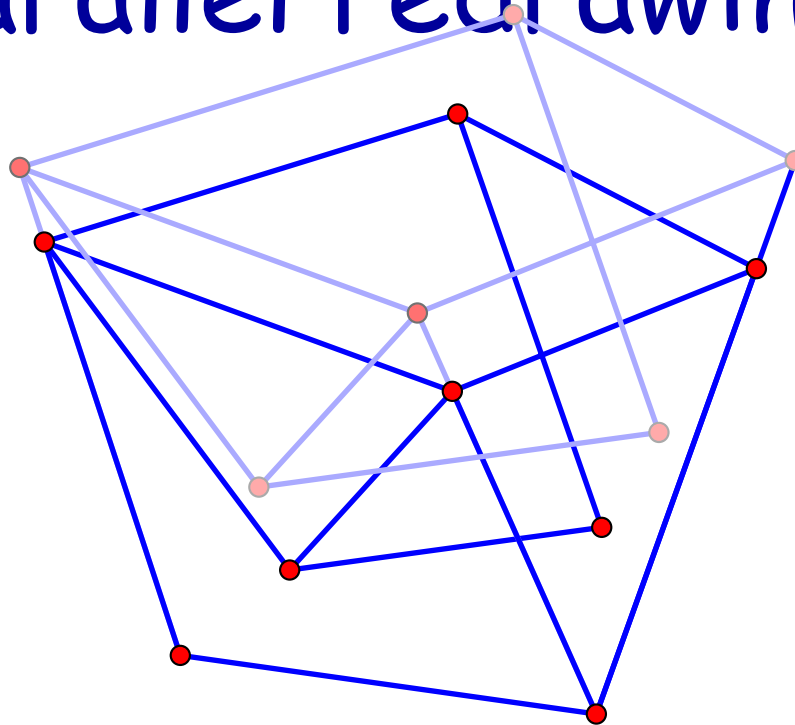
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Direction networks and their parallel redrawings



Parallel redrawing of an embedded graph:

- Another realization of its underlying direction network
- Always can obtain similar ones by translation and rescaling
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Realization space and Configuration space

- **Realization space:** set of all possible realizations of a direction network
- Linear subspace of \mathbb{R}^{2n} : solutions of homogeneous linear system:
 - For all edges ij : $\langle p_i - p_j, d_{ij}^\perp \rangle = 0, \forall ij \in E$
 - $Rp=0$, where R is the "parallel redrawing matrix"

		i		j	
ij	0	d_{ij}^\perp	0	$-d_{ij}^\perp$	0

Two columns per vertex

One row per edge

Realization space and Configuration space

of a **1dof parallel redrawing Laman mechanism**

- **Realization space:** set of all possible realizations
- Linear subspace of \mathbb{R}^{2n} : solutions of homogeneous linear system:
 - For all edges $\langle p_i - p_j, d_{ij}^\perp \rangle = 0, \forall ij \in E$
 - $Rp=0$, where R is the "parallel redrawing matrix"
- **Factor out translations:** pin down a vertex, e.g. $p_1=0$. Still homogeneous, $2n-4$ eqs, $2n-2$ variables.
- **Configuration space: factor out scalings**
 - **Projective view:** factor out scalings (trivial parallel redrawings). Projective line in $2n-3$ dim projective space.
 - **Affine view:** eliminate a "point at infinity". E.g. $x_2=1$. Affine line in Euclidean $2n-3$ space
 - **Oriented-projective view:** factor by positive scalings. Oriented-projective line = great circle on $2n-3$ sphere.

$$R_p=0 \text{ and } R_p=b$$

- $R_p=0$ for a mechanism gives the realization space.
- $R_p=b$ for a Laman graph, b all-but-one-zero vector, and one component acting as "time parameter":
 - Captures an affine part of the configuration space

Generic: R max rank

$$R_p=0 \text{ and } R_p=b$$

- $R_p=0$ for a mechanism gives the realization space.
- $R_p=b$ for a Laman graph, b all-but-one-zero vector, and one component acting as "time parameter":
 - Captures an affine part of the configuration space

Plan

- **Configuration spaces** of parallel redrawing Laman graphs and 1dof Laman mechanisms
- **Parallel Redrawing Sweep** for 1dof Laman mechanisms
- **Pointed pseudo-triangulation** mechanisms
- Further problems: **kinetic point sets** and **graphs, combinatorial sweeps**

The Parallel Redrawing Sweep:

Visualizing the Configuration space of a 1dof Laman mechanism

Configuration space:

- **Projective view:** factor out scalings (trivial parallel redrawings). Projective line in $2n-3$ dim projective space.
- **Affine view:** eliminate a "point at infinity" by pinning down one edge. Affine line in Euclidean $2n-3$ space.

Lemma: Each point traces a linear trajectory projection of config. space on the \mathbb{R}^2 of the 2 coordinates of the point.

Lemma: Points move with constant velocities.

- **Oriented-projective view:** factor by positive scalings. Oriented-projective line = great circle on $2n-3$ sphere. Points trace ellipses

The Parallel Redrawing Sweep:

Visualizing the Configuration space of a 1dof Laman mechanism

Configuration space:

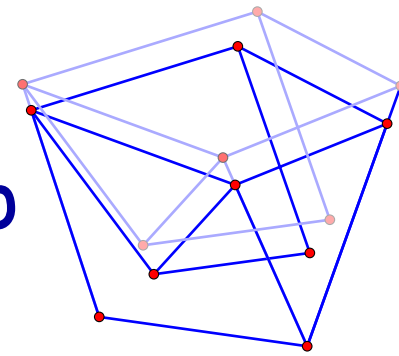
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The parallel redrawing sweep

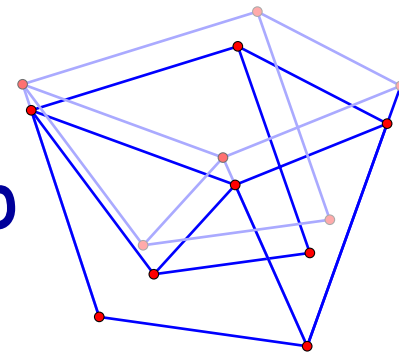


- Sweep of an affine part of the configuration space:
 - Determined by:
 - Choice of a rigid component (to be the "point at infinity"): "pinned down" edge.
 - Choice of an incident edge to "drive the sweep" (time parameter)
 - Sequence of **collision events**: rigid components collapse
- Understanding the sequence of collision events:
 - A rigid component "reverses" (rotates by 180 degrees, i.e. scales by a negative factor)

Lemma: At a collision event, the **contraction** of G on the collapsed edges is a **Laman graph**.

Lemma: The (constant) velocities are the "coordinates" of an "embedding" of the collapsed Laman graph on the r -component at "infinity".

The parallel redrawing sweep

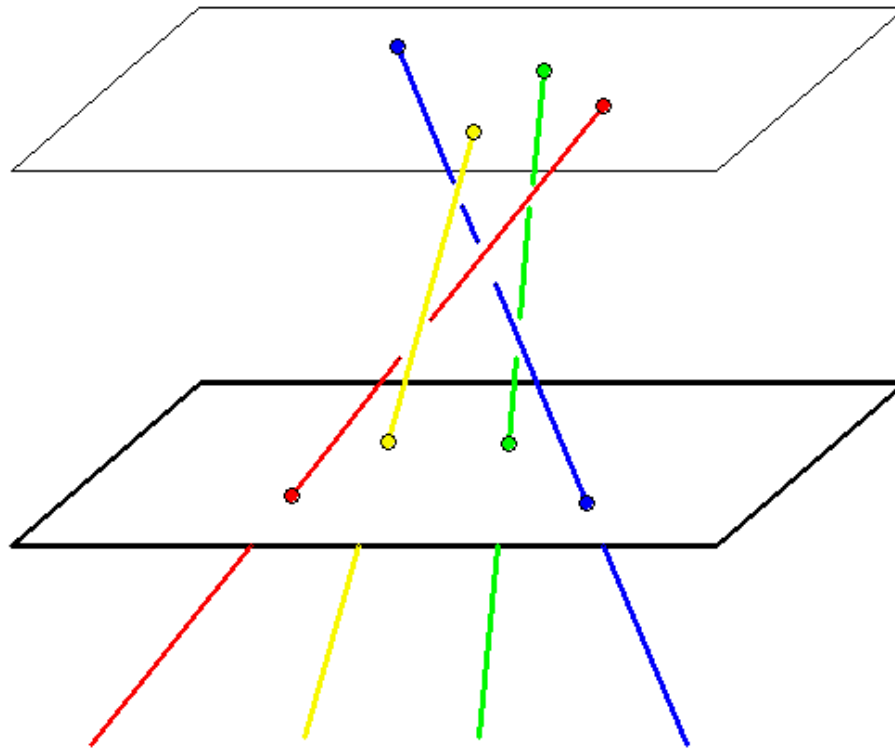


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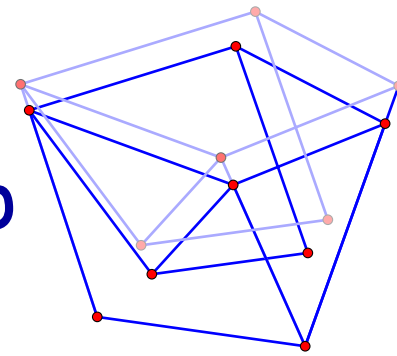
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The parallel redrawing sweep: 3d (space-time) view



A plane sweep of a SPECIAL 3d line arrangement

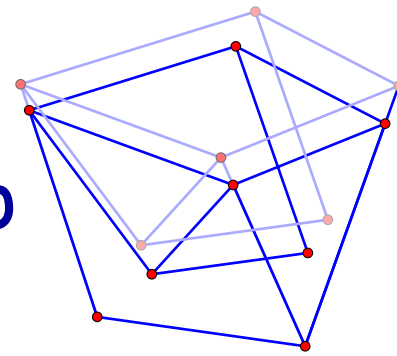
The parallel redrawing sweep



Algorithmic aspects:

- Computing the (combinatorial) events: rigid components of a Laman mechanism
- Predicting the next collision :
 - Linear algebra
 - Can it be done combinatorially?

The parallel redrawing sweep



Algorithmic aspects:

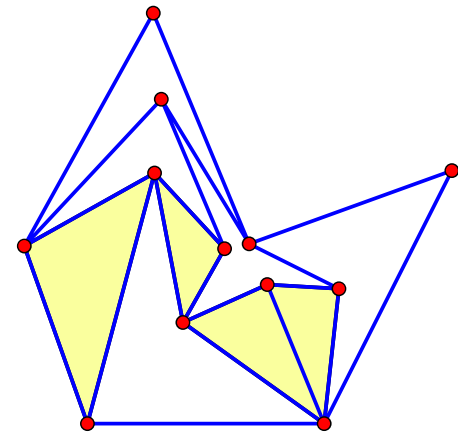
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 - Linear algebra
 - Can it be done combinatorially?

As in "topological sweep" versus "line sweep". **Yes**, for pseudo triangulation mechanisms

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Parallel Redrawing Pointed pseudo-triangulation 1dof-mechanisms



- Pointed Pseudo-triangulation mechanisms (1dof, convex hull missing edge) are **non-crossing throughout a parallel sweep**

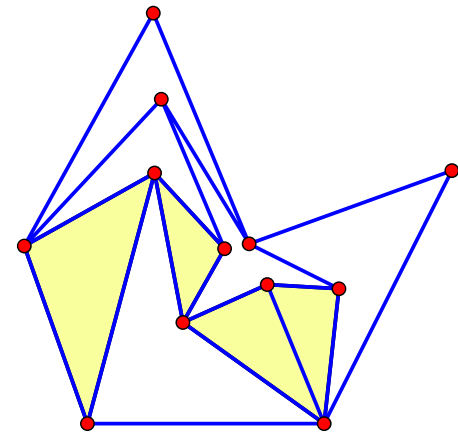
Proof:

- Use *expansiveness* property of pseudo-triangulation mechanisms:

$$\begin{aligned}\langle p_i - p_j, v_i - v_j \rangle &= 0, \quad \forall ij \in E \\ \langle p_i - p_j, v_i - v_j \rangle &\geq 0, \quad \forall ij \notin E\end{aligned}$$

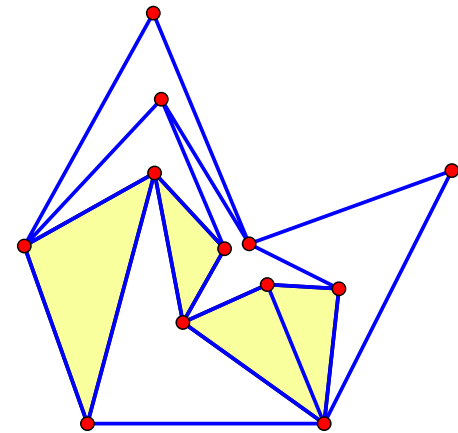
- Interpreted in parallel redrawing setting: **all are pseudo-triangulations**
- Look at oriented matroidal invariants maintained through the parallel sweep (partial unsigned hyperlines / combinatorial pseudo-triangulations): the **facial structure is maintained**

Parallel Redrawing Pointed pseudo-triangulation 1dof-mechanisms



- Lemma: $\sigma_{ij} = \langle p_i - p_j, d_{ij}^\perp \rangle$
do not change during pr-sweep
- Corollary: signs don't change. Hence all positive (expansive). Hence it is a pt-mechanism.
- Lemma: next event only among the "flippable" r-components. Must be extreme at incident joints, and incident components oriented the same way. Property maintained after the event ("flip").
- Lemma: Planar face structure doesn't change. Only combinatorial pseudo-triangulation.
- Obs: Captures partial oriented matroid information (signed hyperlines) of embedded graph.

Parallel Redrawing Pointed pseudo-triangulation 1dof-mechanisms



- Pointed Pseudo-triangulation mechanisms (1dof, convex hull missing edge) are **non-crossing throughout a parallel sweep**
- Replacing a rigid component by another non-crossing graph (on the same kinetic point set) is again a kinetic non-crossing graph

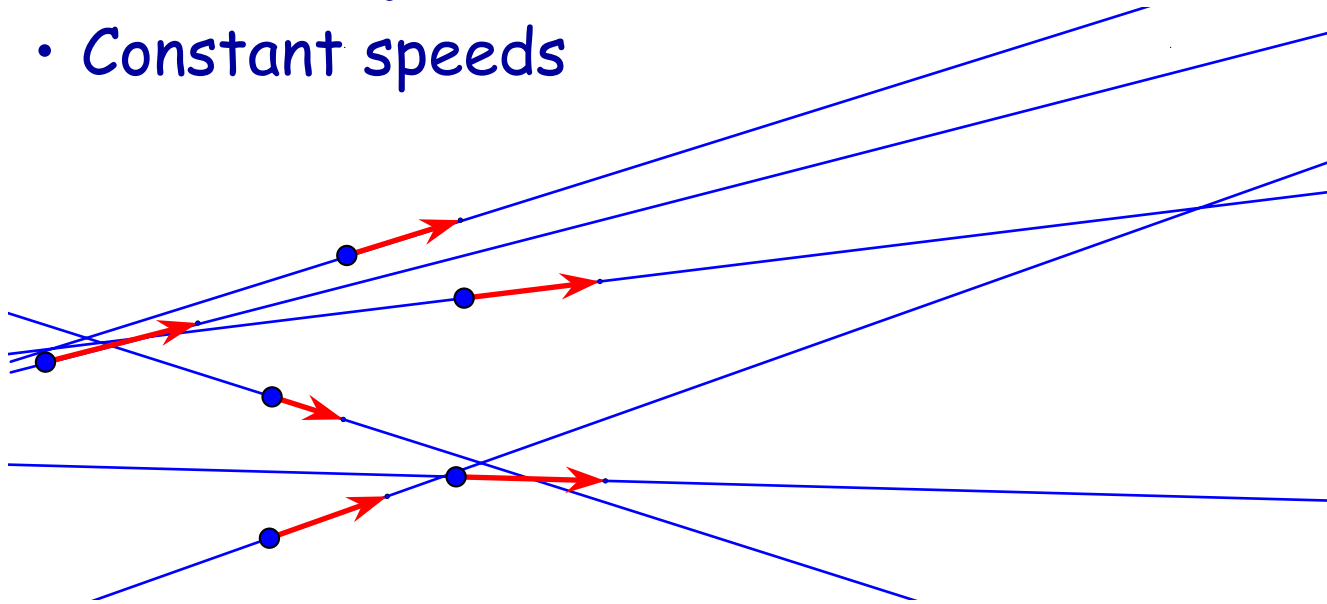
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Study collisions in:

Kinetic Point Sets

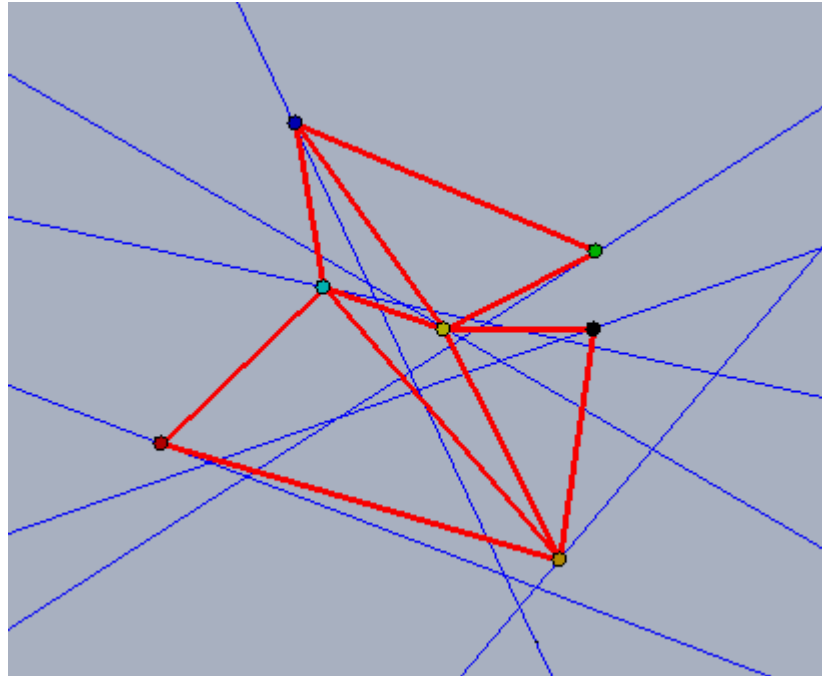
- Points in continuous motion
- Moving with constant velocities
 - Linear trajectories
 - Constant speeds



Study crossings in:

Kinetic Graphs

- Graphs drawn (embedded) on kinetic point sets



Study:

Combinatorial (Topological) Parallel Redrawing Sweep

- For pointed **combinatorial** pseudo-triangulation mechanisms
- Next event predicted combinatorially
- Is every combinatorial sequence realizable?

Questions?

<http://cs.smith.edu/~streinu/Research/KineticPT>